Computing Noise Robustness of Incompatible Quantum Measurements

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MTNS Cambridge, August 22, 2024

Incompatibility in QM

Free spectrahedra

Connecting the two

Computing the noise robustness of incompatibility



Incompatibility in QM

Quantum states and measurements

- Motivation: Classical state \rightsquigarrow probability distributions: $p \in \mathbb{R}^d$, $p \ge 0$, $\sum_i p_i = 1$
- Quantum states \rightsquigarrow density matrices: $\rho \in \mathcal{M}_d(\mathbb{C})$, $\rho \ge 0$, Tr $\rho = 1$
- Measurement outcomes are labeled $\{1, \ldots, k\}$, need to be assigned probabilities
- Measurements: Tuples of matrices (E₁,..., E_k) such that (Tr[E₁ρ],..., Tr[E_kρ]) is a probability distribution for all states ρ
 - $\operatorname{Tr}[E_i \rho] \in \mathbb{R} \rightsquigarrow E_i = E_i^*$
 - $\operatorname{Tr}[E_i \rho] \geq 0 \rightsquigarrow E_i \geq 0$
 - $\sum_{i} \operatorname{Tr}[E_i \rho] = 1 \rightsquigarrow \sum_{i} E_i = I_d$
- Tuples of PSD matrices summing to identity are called positive operator-valued measures (POVMs)
- We call $0 \le E \le I$ quantum effects

Quantum measurements: Compatibility

 Quantum measurements ~>> give the probabilities of the classical outcomes when a quantum state enters a measurement apparatus. Mathematically, measurements are modeled by POVMs

Definition

Two POVMs, $A = (A_1, \ldots, A_k)$ and $B = (B_1, \ldots, B_l)$, are called compatible if there exists a third POVM $C = (C_{ij})_{i \in [k], j \in [l]}$ such that

$$orall i \in [k], \quad A_i = \sum_{j=1}^l C_{ij} \qquad ext{and} \qquad orall j \in [l], \quad B_j = \sum_{i=1}^k C_{ij}$$

The definition generalizes to g-tuples of POVMs $A^{(1)}, \ldots, A^{(g)}$, having respectively k_1, \ldots, k_g outcomes, where the joint POVM C has outcome set $[k_1] \times \cdots \times [k_g]$.

• Other way to say that: jointly measurable

What does it mean?



- Compatible measurements can be simulated by a single joint measurement, by classically post-processing its outputs
- Examples:
 - 1. Trivial POVMs $A = (p_i I_d)$ and $B = (q_j I_d)$ are compatible
 - 2. Commuting POVMs $[A_i, B_j] = 0$ are compatible
 - 3. If the POVM A is projective, then A and B are compatible iff they commute

Noisy POVMs

- POVMs can be made compatible by adding noise, i.e. mixing in trivial POVMs
- Example: dichotomic POVMs and white noise, $s \in [0,1]$

$$(E, I-E) \mapsto s(E, I-E) + (1-s)(\frac{l}{2}, \frac{l}{2})$$
 or $E \mapsto sE + (1-s)\frac{l}{2}$

- Taking s = 1/2 suffices to render any pair of dichotomic POVMs compatible \rightsquigarrow define $C_{ij} := (E_i + F_j)/4$
- $\bullet\,$ For most of the talk, we focus on dichotomic (YES/NO) POVMs

Definition

The incompatibility degree for g measurements on \mathbb{C}^d is the number

$$\gamma(g,d):= \max\{s\in [0,1]\,:\, ext{for all quantum effects } E_1,\ldots,E_g\in \mathcal{M}_d(\mathbb{C}),$$

the noisy versions $sE_i + (1-s)I_d/2$ are compatible}

Free spectrahedra

• A spectrahedron is given by PSD constraints: for

$$egin{aligned} \mathcal{A} &= (\mathcal{A}_1, \dots, \mathcal{A}_g) \in (\mathcal{M}_d(\mathbb{C})^{\mathrm{sa}})^g \ & \mathcal{D}_\mathcal{A}(1) := \left\{ x \in \mathbb{R}^g \ : \ \sum_{i=1}^g x_i \mathcal{A}_i \leq I_d
ight\} \end{aligned}$$



- $\mathcal{D}_{(\sigma_X,\sigma_Y,\sigma_Z)}(1) = \{(x,y,z) \in \mathbb{R}^3 : x\sigma_X + y\sigma_Y + z\sigma_Z \le l_2\} = \text{Bloch ball}$
- A free spectrahedron is the matricization of a spectrahedron

$$\mathcal{D}_A := \bigsqcup_{n=1}^\infty \mathcal{D}_A(n) \quad ext{ with } \quad \mathcal{D}_A(n) := \left\{ X \in (\mathcal{M}_n(\mathbb{C})^{\mathrm{sa}})^g \ : \ \sum_{i=1}^g X_i \otimes A_i \leq I_{nd}
ight\}$$

The matrix diamond is the free spectrahedron defined by

$$\mathcal{D}_{\diamondsuit,g} := \bigsqcup_{n=1}^{\infty} \{ X \in (\mathcal{M}_n(\mathbb{C})^{\mathrm{sa}})^g : \sum_{i=1}^g \epsilon_i X_i \leq I_n, \quad \forall \epsilon \in \{\pm 1\}^g \}$$



- At level one, $\mathcal{D}_{\diamondsuit,g}(1)$ is the unit ball of the ℓ^1 norm on \mathbb{R}^g
- As a free spectrahedron, it is defined by 2^g × 2^g diagonal matrices D_{◊,g} = D_{L1,...,Lg}, with L_i = I₂ ⊗ · · · ⊗ I₂ ⊗ diag(1, −1) ⊗ I₂ ⊗ · · · ⊗ I₂

Spectrahedral inclusion

- Consider two free spectrahedra defined by (A_1, \ldots, A_g) and (B_1, \ldots, B_g)
- We write $\mathcal{D}_A \subseteq \mathcal{D}_B$ if, for all $n \geq 1$, $\mathcal{D}_A(n) \subseteq \mathcal{D}_B(n)$
- Clearly, D_A ⊆ D_B ⇒ D_A(1) ⊆ D_B(1). For the converse implication to hold, one may need to shrink D_A...

Definition

For a free spectrahedron \mathcal{D}_A , we define its inclusion constant as

$$egin{aligned} \delta_{A}(g,d) &:= \max\{s \in [0,1] : ext{for all } g ext{-tuples } B_{1},\ldots,B_{g} \in \mathcal{M}_{d}(\mathbb{C})^{ ext{sa}} \ &\mathcal{D}_{A}(1) \subseteq \mathcal{D}_{B}(1) \implies s.\mathcal{D}_{A} \subseteq \mathcal{D}_{B} \end{aligned}$$

 We shall be concerned with the inclusion constant for the matrix diamond, which we denote by δ(g, d)

Connecting the two

Compatibility in QM \iff matrix diamond inclusion

To a g-tuple $E \in (\mathcal{M}_d(\mathbb{C})^{\operatorname{sa}})^g$, we associate: $\mathcal{D}_{2E-I} := \bigsqcup_{n=1}^{\infty} \{ X \in (\mathcal{M}_n(\mathbb{C})^{\operatorname{sa}})^g : \sum_{i=1}^g X_i \otimes (2E_i - I_d) \leq I_{nd} \}$

Theorem

Let $E \in (\mathcal{M}_d(\mathbb{C})^{sa})^g$ be g-tuple of selfadjoint matrices. Then:

- The matrices E are quantum effects $\iff \mathcal{D}_{\diamondsuit,g}(1) \subseteq \mathcal{D}_{2E-I}(1)$
- The matrices E are compatible quantum effects $\iff \mathcal{D}_{\diamondsuit,g} \subseteq \mathcal{D}_{2E-I}$

At the intermediate levels $1 \le n \le d$, $\mathcal{D}_{\diamondsuit,g}(n) \subseteq \mathcal{D}_{2E-I}(n)$ iff for all isometries $V : \mathbb{C}^n \to \mathbb{C}^d$, the compressed effects $V^* E_i V$ are compatible.

Moreover, the incompatibility degree is equal to the inclusion constant of the matrix diamond: $\forall g, d, \gamma(g, d) = \delta(g, d)$.

Many things are known about the matrix diamond:

- For all $g, d, \tau(d) \leq \delta(g, d), \tau(d) \approx \sqrt{\frac{2}{\pi d}}$ asymptotically (AB and Nechita, 2022)
- For all $g, d, \frac{1}{\sqrt{g}} \leq \delta(g, d)$ (Passer *et al.*, 2018)

Theorem (Passer et al., 2018)

For all g and $d \geq 2^{\lceil (g-1)/2 \rceil}$, $\gamma(g,d) = \delta(g,d) = \frac{1}{\sqrt{g}}$

Phase diagram



- Connection to free spectrahedra also holds for arbitrary outcomes
- Instead of matrix diamond, consider its generalization, the matrix jewel
- The smallest unknown case is d = 2, g = 4 (4 qubit measurements)

Computing the noise robustness of incompatibility

4 qubit measurements: Reminder

- Connection between compatibility and free spectrahedra: The matrices E are compatible quantum effects ⇐⇒ D_{◊,g} ⊆ D_{2E-I}
- $B_i = 2E_i I$ for some $I \ge E_i \ge 0 \iff -I \le B_i \le I$
- $X \in \mathcal{D}_{\diamondsuit,4}(2) \iff \sum_{i=1}^{4} \epsilon_i X_i \le I_2 \quad \forall \epsilon \in \{\pm 1\}^4$
- $X \in \mathcal{D}_{2E-I} \iff \sum_{i=1}^{4} B_i \otimes X_i \le I$

Goal: Maximize

$$\lambda_{\max}\left(\sum_{i=1}^4 B_i\otimes X_i\right)$$

over all allowed X, B as above.

Value: $1/\delta(4,2) = 1/\gamma(4,2)$

4 qubit measurements: Optimization problem

$$\begin{array}{ll} \text{maximize} & \lambda_{\max} \left(\sum_{i=1}^{4} B_i \otimes X_i \right) \\ \text{subject to} & B_1 \leq l_2, \quad -B_1 \leq l_2, \\ & B_2 \leq l_2, \quad -B_2 \leq l_2, \\ & B_3 \leq l_2, \quad -B_3 \leq l_2, \\ & B_4 \leq l_2, \quad -B_4 \leq l_2, \\ & \sum_{i=1}^{4} \epsilon_i X_i \leq l_2, \quad \forall \epsilon \in \{\pm 1\}^4 \\ & B_i, \ X_i \in \mathcal{M}_2(\mathbb{C})^{\text{sa}} \quad \forall i \in \{1, 2, 3, 4\} \end{array}$$

Problem: We need to keep the dimension fixed, otherwise the result is 1/2

• Use Pauli matrices to write

$$B_{i} = b_{0}^{(i)} l_{2} + b_{1}^{(i)} \sigma_{X} + b_{2}^{(i)} \sigma_{Y} + b_{3}^{(i)} \sigma_{Z}$$
$$X_{i} = x_{0}^{(i)} l_{2} + x_{1}^{(i)} \sigma_{X} + x_{2}^{(i)} \sigma_{Y} + x_{3}^{(i)} \sigma_{Z}$$

- We obtain a polynomial optimization problem in commuting variables
- Solve it with Lasserre-Parillo hierarchy

$$\begin{array}{ll} \text{maximize} & \sum_{i} \sum_{j} b_{j}^{(i)} x_{j}^{(i)} \\ \text{subject to} & 1 - (b_{1}^{(0)})^{2} - (b_{2}^{(0)})^{2} - (b_{3}^{(0)})^{2} \geq 0, \\ & (1 - b_{0}^{(i)})^{2} - (b_{1}^{(0)} - b_{1}^{(i)})^{2} - (b_{2}^{(0)} - b_{2}^{(i)})^{2} - (b_{3}^{(0)} - b_{3}^{(i)})^{2} \geq 0, \quad \forall i \in [4] \\ & (1 + b_{0}^{(i)})^{2} - (b_{1}^{(0)} + b_{1}^{(i)})^{2} - (b_{2}^{(0)} + b_{2}^{(i)})^{2} - (b_{3}^{(0)} + b_{3}^{(i)})^{2} \geq 0, \quad \forall i \in [4] \\ & 1 \geq b_{0}^{(i)} \geq -1, \qquad \forall i \in \{1, 2, 3, 4\} \\ & (1 - \sum_{i} \epsilon_{i} x_{0}^{(i)})^{2} \geq \sum_{j} \left(\sum_{i} \epsilon_{i} x_{j}^{(i)}\right)^{2}, \qquad \forall \epsilon \in \{\pm 1\}^{4} \\ & 1 \geq \sum_{i} \epsilon_{i} x_{0}^{(i)}, \qquad \forall \epsilon \in \{\pm 1\}^{4} \\ & b_{i}^{(i)}, \quad x_{i}^{(i)} \in \mathbb{R} \quad \forall i, j \end{array}$$

Conjectured noise robustness of 4 qubit measurements

Known constraints:

- $\gamma(g,d) \geq 1/\sqrt{g} \implies \gamma(4,2) \geq 1/2$
- $\gamma(g, d) \ge \gamma(g+1, d)$ and $\gamma(3, 2) = 1/\sqrt{3}$ $\implies \gamma(4, 2) \le 1/\sqrt{3} \approx 0.58$

Conjecture: $\gamma(4,2) = 2/\sqrt{13} \approx 0.55$

- Theory of extreme points for free spectrahedra: $\gamma(4,2) \leq 2/\sqrt{13}$
- Polynomial optimization: Almost matching lower bounds



- The connection between compatibility and free spectrahedra also holds for more outcomes
- Need to substitute the matrix diamond by the matrix jewel
- Can for example consider one measurement with 2 and one measurement with k outcomes
- Not much is known for this problem
- Can write down a polynomial optimization problem again to compute the noise robustness

Polynomial optimization problem for 2 + 3 outcomes

ma

$$\begin{array}{ll} \text{maximize} & \lambda_{\max} \left(\sum_{i=1}^{3} B_i \otimes X_i \right) \\ \text{subject to} & B_1 \leq I_m, \quad -B_1 \leq I_m, \\ & -\frac{3}{2} B_2 \leq I_m, \quad -\frac{3}{2} B_3 \leq I_m, \\ & \frac{3}{2} B_2 + \frac{3}{2} B_3 \leq I_m, \\ & \pm X_1 + \frac{4}{3} X_2 - \frac{2}{3} X_3 \leq I_m, \\ & \pm X_1 - \frac{2}{3} X_2 + \frac{4}{3} X_3 \leq I_m, \\ & \pm X_1 - \frac{2}{3} X_2 - \frac{2}{3} X_3 \leq I_m, \\ & B_i, \ X_i \in \mathcal{M}_m(\mathbb{C})^{\text{sa}} \quad \forall i \in \{1, 1\} \end{array}$$

If we don't fix the dimension m, we can solve this using the NPA hierarchy or the hierarchy for solving trace polynomials by Klep et al., 2022

2, 3

Conjectured noise robustness of 2 + k outcomes measurements

• We know that
$$1/2 \leq \gamma((2,3),m) \leq 1/\sqrt{2}$$

Conjecture:
$$\gamma((2,k),m)=rac{1}{2}\left(1+rac{1}{1+\sqrt{k}}
ight)$$

For $k \in \{2, 3, 4, 5\}$:

- Theory of extreme points for free spectrahedra: $\gamma((2,k),m) \leq \frac{1}{2} \left(1 + \frac{1}{1+\sqrt{k}}\right)$
- Corresponding measurements: $\{|0\rangle\!\langle 0|\,,|1\rangle\!\langle 1|\}$ and $\{|+\rangle\!\langle +|\,,|-\rangle\!\langle -|\,,0,\dots,0\}$
- NPA hierarchy almost matching lower bounds up to high precision

- Can we simplify the optimization problems further using knowledge about extreme points of free spectrahedra?
- Can we extract exact certificates from the polynomial optimization?
- Can we solve the polynomial optimization problem more efficiently (allowing to tackle larger problems)?

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