Matrix convex sets in quantum information theory: Measurement compatibility and beyond

Andreas Bluhm Univ. Grenoble Alpes, CNRS, Grenoble INP, LIG

February 14, 2023, Toulouse



Matrix convex sets

We consider free sets:

$$\mathcal{F}=\bigsqcup_{i\in\mathbb{N}}\mathcal{F}_i,$$

where $\mathcal{F}_i \subseteq (\mathcal{M}_i^{\mathrm{sa}})^g$.

The free set \mathcal{F} is matrix convex if it is closed under direct sums and unital completely positive maps:

- $(A_1,\ldots,A_g) \in \mathcal{F}_i, (B_1,\ldots,B_g) \in \mathcal{F}_j \implies (A_1 \oplus B_1,\ldots,A_g \oplus B_g) \in \mathcal{F}_{i+j}.$
- $(A_1, \ldots, A_g) \in \mathcal{F}_i, \ \Phi : \mathcal{M}_i \to \mathcal{M}_j \ \mathsf{UCP} \implies (\Phi(A_1), \ldots, \Phi(A_g)) \in \mathcal{F}_j$

UCP maps $\Phi : \mathcal{M}_d \to \mathcal{M}_m$ are maps such that $\Phi \otimes \mathrm{id}_n$ is positive for all $n \in \mathbb{N}$ and $\Phi(I_d) = I_m$.

Alternatively, $\Phi(X) = \sum_i K_i^* X K_i$ such that $\sum_i K_i^* K_i = I_m$, $K_i \in \mathcal{M}_{d,m}$.

Minimal and maximal matrix convex sets

- Unless \mathcal{F}_1 is a simplex, there are arbitrarily many different matrix convex sets with the same \mathcal{F}_1 . However, there is a largest and a smallest such set:
- For a closed convex set \mathcal{C} ,

$$\mathcal{W}^{\sf max}_n(\mathcal{C}) := \Big\{ X \in (\mathcal{M}^{
m sa}_n)^{\sf g} : \sum_{i=1}^{\sf g} c_i X_i \leq lpha I \; orall(lpha, {\sf c}) \; {
m supp. hyperplanes for } \mathcal{C} \Big\}$$

 $\bullet\,$ For a closed convex set $\mathcal{C},$

$$\mathcal{W}_n^{\min}(\mathcal{C}) := \Big\{ \sum_j X = z_j \otimes Q_j \in (\mathcal{M}_n^{\operatorname{sa}})^g : z_j \in \mathcal{C}, \ Q_j \ge 0 \ \forall j, \sum_j Q_j = I_n \Big\}$$

Observe \$\mathcal{W}_1^{max}(\mathcal{C}) = \mathcal{C} = \mathcal{W}_1^{min}(\mathcal{C})\$. \$\mathcal{W}^{max}(\mathcal{C})\$ quantizes hyperplanes, \$\mathcal{W}^{min}(\mathcal{C})\$ quantizes extreme points.

Definition

Let $d, g \in \mathbb{N}$ and $\mathcal{C} \subset \mathbb{R}^g$ closed convex. The inclusion set is defined as $\Delta_{\mathcal{C}}(d) := \left\{ s \in [0, 1]^g : s \cdot \mathcal{W}_d^{\mathsf{max}}(\mathcal{C}) \subseteq \mathcal{W}_d^{\mathsf{min}}(\mathcal{C}) \right\}.$ If \mathcal{C} is the ℓ_{∞}^g unit ball, we write $\Delta_{\Box}(g, d)$.

Depending on the set \mathcal{C} , sometimes bounds on the inclusion set are known.

Measurement compatibility

- Motivation: Classical state \rightsquigarrow probability distributions: $p \in \mathbb{R}^d$, $p \ge 0$, $\sum_i p_i = 1$.
- Quantum states \rightsquigarrow density matrices: $\rho \in \mathcal{M}_d(\mathbb{C})$, $\rho \ge 0$, $\operatorname{Tr} \rho = 1$.
- Measurement outcomes are labeled $\{1, \ldots, k\}$, need to be assigned probabilities.
- Measurements: Tuples of matrices (E₁,..., E_k) such that (Tr[E₁ρ],..., Tr[E_kρ]) is a probability distribution for all states ρ.
 - $\operatorname{Tr}[E_i \rho] \in \mathbb{R} \rightsquigarrow E_i = E_i^*$.
 - $\operatorname{Tr}[E_i \rho] \geq 0 \rightsquigarrow E_i \geq 0.$
 - $\sum_{i} \operatorname{Tr}[E_i \rho] = 1 \rightsquigarrow \sum_{i} E_i = I_d.$
- Tuples of PSD matrices summing to identity are called positive operator-valued measures (POVMs).

Quantum measurements: Compatibility

 Quantum measurements ~>> give the probabilities of the classical outcomes when a quantum state enters a measurement apparatus. Mathematically, measurements are modeled by POVMs.

Definition

Two POVMs, $A = (A_1, ..., A_k)$ and $B = (B_1, ..., B_l)$, are called compatible if there exists a third POVM $C = (C_{ij})_{i \in [k], j \in [l]}$ such that

$$\forall i \in [k], \quad A_i = \sum_{j=1}^l C_{ij} \quad \text{and} \quad \forall j \in [l], \quad B_j = \sum_{i=1}^k C_{ij}.$$

The definition generalizes to g-tuples of POVMs $A^{(1)}, \ldots, A^{(g)}$, having respectively k_1, \ldots, k_g outcomes, where the joint POVM C has outcome set $[k_1] \times \cdots \times [k_g]$.

• Other way to say that: jointly measurable.

What does it mean?



- Compatible measurements can be simulated by a single joint measurement, by classically post-processing its outputs $A_i^{(j)} = \sum_{\lambda} p_j(i|\lambda) C_{\lambda}$.
- Examples:
 - 1. Trivial POVMs $A = (p_i I_d)$ and $B = (q_j I_d)$ are compatible.
 - 2. Commuting POVMs $[A_i, B_j] = 0$ are compatible.
 - 3. If the POVM A is projective, then A and B are compatible if and only if they commute.

Noisy POVMs

- POVMs can be made compatible by adding noise, i.e. mixing in trivial POVMs.
- Example: dichotomic POVMs and white noise, $s \in [0, 1]$:

$$E, I-E) \mapsto s(E, I-E) + (1-s)(\frac{l}{2}, \frac{l}{2})$$
 or $E \mapsto sE + (1-s)\frac{l}{2}$.

- Taking s = 1/2 suffices to render any pair of dichotomic POVMs compatible →→ define C_{ij} := (E_i + F_j)/4.
- From now on, we focus on dichotomic (YES/NO) POVMs.

Definition

The compatibility region for g measurements on \mathbb{C}^d is the set

$$\Gamma(g,d):=\{s\in [0,1]^g\,:\, {
m for all quantum effects} \,\, E_1,\ldots,E_g\in \mathcal{M}_d(\mathbb{C})$$

the noisy versions $s_i E_i + (1 - s_i) I_d / 2$ are compatible}

$$\begin{split} \Gamma(g,d) &:= \{s \in [0,1]^g \ : \ \text{for all quantum effects} \ E_1,\ldots,E_g \in \mathcal{M}_d(\mathbb{C}), \\ & \text{the noisy versions} \ s_iE_i + (1-s_i)I_d/2 \ \text{are compatible} \} \end{split}$$

- The set $\Gamma(g, d)$ is convex.
- For all i ∈ [g], e_i ∈ Γ(g, d): every measurement is compatible with g − 1 trivial measurements.
- For d ≥ 2, (1,1,...,1) ∉ Γ(g,d): there exist incompatible measurements.
- For all $d \ge 2$, $\Gamma(2, d)$ is a quarter-circle.



Generally speaking, the set $\Gamma(g, d)$ tells us how robust (to noise) is the incompatibility of g dichotomic measurements on \mathbb{C}^d .

Link measurement compatibility and matrix convex sets

From now on, we concentrate on measurements with two outcomes and identify $E^{(i)} = \{E_i, I - E_i\}$ with E_i .

Theorem

Let

$$A=\sum_{j=1}^g e_j\otimes (2E_j-I).$$

Then,

- 1. $A \in W^{\max}_d(\mathcal{B}(\ell_{\infty}^g))$ if and only if $\{E_j\}_{j \in [g]}$ is a collection of POVMs.
- 2. $A \in W^{\min}_d(\mathcal{B}(\ell_{\infty}^g))$ if and only if $\{E_j\}_{j \in [g]}$ is a collection of compatible POVMs.

Proof sketch

- $\mathcal{W}_d^{\max}(\mathcal{B}(\ell_\infty^g))$ is given in terms of hyperplanes. Have to verify $-I \leq A_i = 2E_i I \leq I$ $\implies 0 \leq E_i \leq I$.
- Reminder:

$$\mathcal{W}^{\min}_n(\mathcal{B}(\ell^g_\infty)) := \Big\{ X = \sum_j z_j \otimes Q_j \in (\mathcal{M}^{\mathrm{sa}}_n)^g : z_j \in \mathcal{C} \,\, \forall j, \,\, Q \,\, \mathrm{POVM} \Big\}.$$

• Going to extreme points:

$$2E_j - I = \sum_{\epsilon \in \{\pm 1\}} \epsilon(j) Q_\epsilon$$

• Using $\sum_{\epsilon} Q_{\epsilon} = I$:

$$E_j = \sum_{\epsilon \in \{\pm 1\}} \delta_{\epsilon(j),1} Q_{\epsilon}.$$

• $\{Q_{\epsilon}\}_{\epsilon}$ is a joint POVM.

Inclusion sets and compatibility regions

Theorem

Let g, $d \in \mathbb{N}$. Let $s \in [0, 1]^g$. Then, $\{s_i E_i + (1 - s_i)I/2\}_{i \in [g]}$ is a collection of compatible POVMs for all POVMs $\{E_i\}_{i \in [g]}$, if and only if $s \in \Delta_{\Box}(g, d)$. An equivalent way to phrase this is $\Gamma(g, d) = \Delta_{\Box}(g, d)$.

• This follows from the computation

$$A'_i = 2(s_i E_i + (1 - s_i)I/2) - I = s_i(2E_i - I) = s_iA_i.$$

- So adding noise means scaling the tensor A and hence s · W^{max}_d(B(l^g_∞)) is the set of noisy measurements.
- Thus, $s \cdot A \in \mathcal{W}^{\min}_d(\mathcal{B}(\ell^g_\infty))$ means the noisy measurements are compatible.

Polytope compatibility

Work in progress with Ion Nechita and Simon Schmidt

Definition

Let \mathcal{P} be a polytope in \mathbb{R}^g such that $0 \in \operatorname{int} \mathcal{P}$. Let

$$A=(A_1,\ldots,A_g)\in \mathcal{M}^{\mathrm{sa}}_d(\mathbb{C})^g\cong \mathbb{R}^g\otimes \mathcal{M}^{sa}_d(\mathbb{C})$$

a g-tuple of Hermitian matrices. Then, A are \mathcal{P} -operators if and only if $A \in \mathcal{W}_d^{\max}(\mathcal{P})$. Moreover, A are \mathcal{P} -compatible if and only if $A \in \mathcal{W}_d^{\min}(\mathcal{P})$.

Motivation:

- A are $\mathcal{B}(\ell_{\infty}^{g})$ -operators if and only if $\frac{1}{2}(A_{i}+I)$ are dichotomic POVMs.
- A are $\mathcal{B}(\ell_{\infty}^{g})$ -compatible if and only if $\frac{1}{2}(A_{i}+I)$ are compatible dichotomic POVMs.

Interlude: General Probabilistic Theories



A GPT is a triple (V, V⁺, 1), where V is a vector space, V⁺ ⊆ V is a cone, and 1 is a linear form on V; A = V^{*}, A⁺ = (V⁺)^{*}, and 1 ∈ A⁺

• The set of states $K := V^+ \cap \mathbb{1}^{-1}(\{1\})$

Equivalent formulation

Theorem

Let d, g, $k \in \mathbb{N}$ and let \mathcal{P} be a polytope with k extremal points $v_1, \ldots, v_k \in \mathbb{R}^g$ such that $0 \in int \mathcal{P}$. Let $A = (A_1, \ldots, A_g) \in \mathcal{M}_d^{sa}(\mathbb{C})^g$ be a g-tuple of Hermitian matrices. Let us consider the map $\mathcal{A} : \mathcal{M}_d^{sa} \to \mathbb{R}^g$,

$$\mathcal{A}(X) = (\operatorname{Tr}[A_1X], \ldots, \operatorname{Tr}[A_gX]).$$

Then,

1. A are \mathcal{P} -operators if and only if \mathcal{A} is a channel between $(\mathcal{M}_d^{sa}, PSD_d, Tr)$ and $(\mathcal{V}(\mathcal{P}), \mathcal{V}(\mathcal{P})^+, \mathbb{1}_{\mathcal{P}})$.

2. A are \mathcal{P} -compatible if and only if in addition \mathcal{A} factors through the k-simplex Δ_k .

Interpretation: \mathcal{P} defines some kind of allowed post-processing. In the case of $\mathcal{B}(\ell_{\infty}^{g})$: Classical post-processing.

Magic squares

A magic square is a collection of positive operators A_{ij} , $i, j \in [N]$, such that

The magic square is said to be semiclassical if

$$A = \sum_{i,j \in [N]} |i\rangle \langle j| \otimes A_{ij} = \sum_{\pi \in \mathcal{S}_N} P_{\pi} \otimes Q_{\pi},$$

where P_{π} is the permutation matrix associated to π and $\{Q_{\pi}\}_{\pi}$ is a POVM.

Birkhoff polytope compatibility

Definition

For a given $N \ge 2$, the Birkhoff body $\mathcal{B}_N(1)$ is defined as the set of $(N-1) \times (N-1)$ truncations of $N \times N$ bistochastic matrices, shifted by J/N:

$$\mathcal{B}_{N} = \{ \mathcal{A}^{(N-1)} - J_{N-1}/N \, : \, \mathcal{A} \in \mathcal{M}_{N}(\mathbb{R}) \, \text{ bistochastic} \} \subset \mathbb{R}^{(N-1)^{2}}$$

Theorem

Consider a $(N-1)^2$ -tuple of selfadjoint matrices $A \in \mathcal{M}_d^{sa}(\mathbb{C})^{(N-1)^2}$ and the corresponding matrix $\tilde{A} \in \mathcal{M}_N(\mathcal{M}_d(\mathbb{C}))$. Then:

1. The matrix \tilde{A} is a magic square if and only if A - I/N are \mathcal{B}_N -operators.

2. The matrix \tilde{A} is a semiclassical magic square if and only if A - I/N are \mathcal{B}_N -compatible.

Relation to measurement incompatibility

Is being a semiclassical magic square the same as being compatible? No.



These measurements are compatible, but they do not form a semiclassical magic square.

Reason: \mathcal{B}_N -compatibility restricts the post-processing to $p_i(j|\lambda) = p_j(i|\lambda)$, i.e., enforces special structure in the joint POVM.

Measurement compatibility with shared effects

Can we generalize the magic square example?

 $\mathcal{P} = (-1/3, -1/3, -1/3) + \operatorname{conv}\{((1, 0, 0), (0, 0, 0), (0, 0, 1), (0, 1, 0), (0, 1, 1)\}.$



Consider $(A, B, C) \in (\mathcal{M}_d(\mathbb{C})^{sa})^3$. Then, we have $(A, B, C) + 1/3(I, I, I) \in \mathcal{W}_d^{\max}(\mathcal{P})$ if and only if both $(A, B, I_d - A - B)$ and $(A, C, I_d - A - C)$ are POVMs.

When does $(A, B, C) + 1/3(I, I, I) \in W_d^{\min}(\mathcal{P})$ hold? Equivalent to the existence of a joint measurement such that

$$\begin{array}{c|cccc} Q_1 & 0 & 0 \\ \hline Q_1 & Q_5 & Q_4 \\ \hline 0 & Q_5 & Q_4 \\ \hline 0 & Q_3 & Q_2 \\ \hline = A & = C & = I_d - A - C \end{array} = A$$

Not all joint measurements are of this form, check

$$\left(\frac{1}{2}\textit{I}_2,\frac{1}{2}\ket{0}\!\!\bra{0},\frac{1}{2}\ket{1}\!\!\bra{1}\right) \qquad \text{ and } \qquad \left(\frac{1}{2}\textit{I}_2,\frac{1}{2}\ket{+}\!\!\bra{+},\frac{1}{2}\ket{-}\!\!\bra{-}\right).$$

- Measurement incompatibility can be phrased as inclusion of matrix convex sets. Base set: cube.
- Noise robustness corresponds to inclusion constants.
- Generalization: \mathcal{P} -operators and \mathcal{P} -compatible operators.
- Examples include magic squares and compatibility with shared elements (under restricted post-processing).

Can we find more tasks in quantum information theory which can be formulated as $\mathcal{P}\text{-}compatibility?}$